100

SENDER 102S

102-2
CLIENT

102-1
CLIENT

104-1
108-1
108-2
CHECK
SERVER

CHECK
SERVER 104-2

106

108-6 104-6
CHECK
SERVER

104-5
104-3
108-4
CHECK
SERVER 104-4

SERVER

SERVER

NETWORK

102-4
CLIENT

102-3
CLIENT

RECIPIENT 102R

FIG. 1

270

200

PROCESSOR

240

DISK

220

MEMORY

260

NETWORK
INTERFACE

$F_{IG,}$ 2

GENERATE A MESSAGE M TO BE ENCRYPTED — 300

PICK A RANDOM ELEMENT $k \in [0 \ldots q-1]$ — 302

COMPUTE A SYMMETRIC KEY $K = HASH(g^k \mod p)$ SUCH THAT K IS A VALID ENCRYPTION KEY FOR A GIVEN ENCRYPTION TECHNIQUE — 304

COMPUTE A QUINTUPLE (a, b, M', c, C_d) WHERE $a = y_d^{\alpha} g^k$, $b = g^{\alpha}$, $M' = E_K(M)$, c IS A PROOF OF KNOWLEDGE OF (α, k), AND C_d IS A CERTIFICATE ON PUBLIC KEY y_d — 306

SEND (a, b, M', c, C_d) THROUGH THE NETWORK WITH INFORMATION IDENTIFYING SENDER AND RECIPIENT — 308
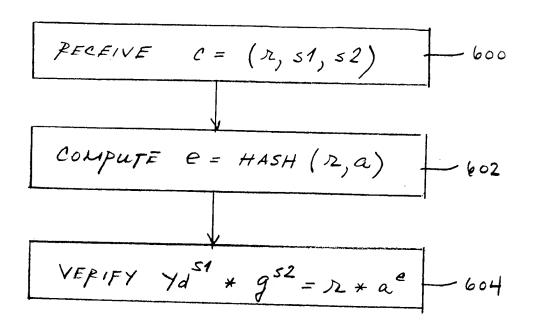
FIG. 3

RECEIVE (a, b, M') FROM SERVER ——— 400

COMPUTE $B = b^{x_d} \pmod{p}$ ——— 402

COMPUTE $K = $ HASH $(a/B \bmod p)$ ——— 404

COMPUTE $M = D_K(M')$ ——— 406

OUTPUT PLAINTEXT MESSAGE M ——— 408

$FIG. 4$

SELECT $\beta 1, \beta 2 \in [0 \ldots q-1]$ — 500

COMPUTE $r = Y_d{}^{\beta 1} g^{\beta 2} \pmod{p}$ — 502

COMPUTE $e = HASH(r, a)$ — 504

COMPUTE $s1 = \beta 1 + e * \alpha \pmod{q}$ — 506

COMPUTE $s2 = \beta 2 + e * k \pmod{q}$ — 508

OUTPUT $(r, s1, s2)$ AS PROOF C — 510

FIG. 5

RECEIVE $C = (r, s1, s2)$ —— 600

COMPUTE $e = HASH(r, a)$ —— 602

VERIFY $yd^{s1} * g^{s2} = r * a^{e}$ —— 604

FIG. 6